

# Lecture 5: Blowups

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The goal of this course will be derived blowups. In this lecture we discuss the 1-categorical version. At the end we paint a picture of the deformation to the normal cone construction.

## 1 Blowups of affine schemes

**Definition 1.** Suppose that  $R$  is a ring and  $I \subseteq R$  an ideal. Similar to  $\mathbb{P}^n$ , we define a presheaf on  $\mathcal{A}ff/R$  by sending a ring homomorphism  $R \rightarrow A$  to the set

$$Bl_R I(A) := \{A \otimes_R I \twoheadrightarrow L\} / \sim$$

of equivalence classes of surjections  $A \otimes_R I \twoheadrightarrow L$  towards an invertible  $A$ -module  $L$  assuch that: For each  $d$  the morphism induced by  $I \rightarrow A \otimes_R I \rightarrow L$  factors as<sup>1</sup>

$$\begin{array}{ccc} \text{Sym}_R^d(I) & \xrightarrow{\quad} & \text{Sym}_A^d(L) \\ & \searrow \text{Id} \text{---} & \nearrow \end{array}$$

As with  $\mathbb{P}^n$ , one declares two surjections to be equivalent if there exists an isomorphism  $L \xrightarrow{\sim} L'$  of  $A$ -modules making a commutative triangle

$$\begin{array}{ccc} & & L \\ & \nearrow & \downarrow \cong \\ A \otimes_R I & & L' \\ & \searrow & \end{array}$$

**Proposition 2.** For any set of generators  $I = \langle r_\lambda \rangle_{\lambda \in \Lambda}$  there is a closed immersion  $Bl_R I \hookrightarrow \mathbb{P}_R^\Lambda$ . Consequently, the presheaf  $Bl_R I$  is a scheme.

*Proof.* Let  $R^{\oplus \Lambda} \twoheadrightarrow I$  be the surjection of  $R$ -modules induced by the generators  $r_\lambda$ . Then for each  $R \rightarrow A$ , the pullback  $A^{\oplus \Lambda} \rightarrow A \otimes_R I$  is surjective, and sending  $[A \otimes_R I \twoheadrightarrow L]$  to  $[A^{\oplus \Lambda} \twoheadrightarrow A \otimes_R I \twoheadrightarrow L]$  defines a morphism of presheaves

$$Bl_R I \rightarrow \mathbb{P}_R^\Lambda.$$

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<sup>1</sup> Given an  $A$ -module  $M$ , one defines

$$\text{Sym}_A^d M := \text{colim}_{\sigma \in \Sigma_d} \underbrace{M \otimes_A \cdots \otimes_A M}_{d \text{ factors}}$$

where the colimit is over elements in the symmetric group  $\Sigma_d$ , which acts by permuting the factors of  $M^{\otimes d}$ .

We claim that this is a closed immersion. First note that it is a monomorphism, since the outer triangle below is commutative if and only if the inner one is by surjectivity of  $(*)$ .

$$\begin{array}{ccc}
 & & L \\
 & \nearrow & \downarrow \\
 A^{\oplus \Lambda} & \xrightarrow{(*)} A \otimes_R I & \downarrow \cong \\
 & \searrow & L'
 \end{array}$$

Let  $j : \text{Spec}(A) \rightarrow \mathbb{P}^\Lambda$  be any morphism from an affine, with corresponding epimorphism  $[A^{\oplus \Lambda} \twoheadrightarrow L]$ . We wish to find an isomorphism  $\text{Spec}(A/J) \xrightarrow{\sim} \text{Bl}_R I \times_{\mathbb{P}^\Lambda} \text{Spec}(A)$  for some ideal  $J$ . Since  $L$  is an invertible module, it is projective, so there exists an isomorphism  $L \oplus L' \cong A^{\oplus M}$  for some  $A$ -module  $L'$  and set  $M$ . On the other hand, choose a surjection from a free module to the kernel of  $R^{\oplus \Lambda} \twoheadrightarrow I$  so we get an exact sequence

$$R^{\oplus \Lambda'} \rightarrow R^{\oplus \Lambda} \rightarrow I \rightarrow 0.$$

Finally, consider the composition  $\Phi : A^{\oplus \Lambda'} \rightarrow A^{\oplus \Lambda} \twoheadrightarrow L \hookrightarrow A^{\oplus M}$ . The  $A$ -module morphism  $\Phi$  is represented by a matrix  $[\Phi]$ . Let  $J \subseteq A$  be the ideal generated by the coefficients of  $[\Phi]$ . Since  $(A/J) \otimes_A \Phi = 0$  (by definition of  $J$ ) we get a factorisation

$$\begin{array}{ccccccc}
 (A/J)^{\oplus \Lambda'} & \longrightarrow & (A/J)^{\oplus \Lambda} & \longrightarrow & (A/J) \otimes_A I & \longrightarrow & 0 \\
 & \searrow & \parallel & & \downarrow & & \\
 & & (A/J)^{\oplus \Lambda} & \longrightarrow & (A/J) \otimes_A L & \subseteq & (A/J)^{\oplus M} \cong (A/J) \otimes_A (L \oplus L') \\
 & & & & & & \text{(1)}
 \end{array}$$

The surjection  $(A/J) \otimes_A I \twoheadrightarrow (A/J) \otimes_A L$  corresponds to a morphism  $\text{Spec}(A/J) \rightarrow \text{Bl}_R I$ . Moreover, the fact that the composition  $(A/J)^{\oplus \Lambda} \twoheadrightarrow (A/J) \otimes_A I \twoheadrightarrow (A/J) \otimes_A L$  is the pullback of the original  $A^{\oplus \Lambda} \twoheadrightarrow L$  corresponds to the following square being commutative.

$$\begin{array}{ccc}
 \text{Spec}(A/J) & \longrightarrow & \text{Spec}(A) \\
 \downarrow & & \downarrow \\
 \text{Bl}_R I & \longrightarrow & \mathbb{P}_R^\Lambda
 \end{array}$$

So we have found a morphism

$$\text{Spec}(A/J) \rightarrow \text{Bl}_R I \times_{\mathbb{P}_R^\Lambda} \text{Spec}(A).$$

We claim it is an isomorphism. It is clearly injective since  $\text{Spec}(A/J) \rightarrow \text{Spec}(A)$  is injective, so it suffices to show that it is surjective. But this is almost by design. Suppose

$$s : \text{Spec}(B) \rightarrow \text{Bl}_R I \times_{\mathbb{P}^\Lambda} \text{Spec}(A)$$

is any morphism. Since  $\text{Spec}(A/J) \rightarrow \text{Bl}_R I \times_{\mathbb{P}^\Lambda} \text{Spec}(A) \rightarrow \text{Spec}(A)$  are both monomorphisms, to show that  $s$  factors through  $\text{Spec}(A/J)$ , it suffices to show that

$\text{Spec}(B) \rightarrow \text{Spec}(A)$  factors through  $\text{Spec}(A/J)$ . I.e., that the corresponding  $A \rightarrow B$  sends all elements of  $J$  to zero. The original  $s$  corresponds to a commutative square

$$\begin{array}{ccc} \text{Spec}(B) & \longrightarrow & \text{Spec}(A) \\ \downarrow & & \downarrow \\ Bl_R I & \longrightarrow & \mathbb{P}^\Lambda \end{array}$$

The existence of such a square implies there is a factorisation  $B^{\oplus \Lambda} \twoheadrightarrow B \otimes_R I \dashrightarrow B \otimes_A L$ . This implies  $B^{\oplus \Lambda'} \rightarrow B^{\oplus M}$  is zero (cf. the diagram Eq.(1)), which implies all coefficients of  $[\Phi]$  are zero, which implies all elements of  $J$  are sent to zero in  $B$ .  $\square$

**Example 3.**

1. Suppose that the canonical morphism  $A \otimes_R I \rightarrow A$  is an isomorphism (e.g.,  $A = R[r^{-1}]$ ) for some  $r \in I$ ). In this case we are considering surjections  $A \twoheadrightarrow L$ . But any surjection of invertible modules is an isomorphism<sup>2</sup> so in this case  $Bl_R I(A)$  consists of a single element  $[A \otimes_R I \xrightarrow{\sim} A]$ . In particular, every  $r \in I$ , induces an isomorphism

$$U \times_X Bl_R I \xrightarrow{\sim} U$$

where  $U \rightarrow X$  is the open immersion  $j \text{Spec}(R[r^{-1}]) \rightarrow j \text{Spec}(R)$ . In other words,  $Bl_R I \rightarrow j \text{Spec}(R)$  is an isomorphism outside of the closed subscheme  $j \text{Spec}(R/I) \subseteq j \text{Spec}(R)$ .

2. Next, consider the case  $A = R/I$ . In this case  $A \otimes_R I \cong I/I^2$ . If furthermore there are  $t_1, \dots, t_c \in I$  that induce an isomorphism  $I/I^2 \cong (R/I)^{\oplus c}$ , then  $Bl_R I(R/I)$  is the set of equivalence classes of quotients of  $A \otimes_R I \cong I/I^2 \cong (R/I)^{\oplus c}$ . That is,

$$Bl_R I(R/I) \cong \mathbb{P}^{c-1}(R/I).$$

It follows that, in this case, there is an isomorphism

$$Z \times_X Bl_R I \xrightarrow{\sim} Z \times \mathbb{P}^{c-1}$$

where  $Z = j \text{Spec}(R/I)$ .

The geometric interpretation is as follows. The  $R/I$ -module  $I/I^2$  is the module of functions on  $\text{Spec}(R)$  which vanish on  $R/I$ , modulo the relation:  $f \sim g$  if their “linear terms” agree. This is called the *conormal* module. Then a point  $x$  of  $Bl_R I$  over a point  $z$  in  $Z$  is a “normal direction” to  $Z$  at  $z$  up to scalar.

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<sup>2</sup>A homomorphism  $A \rightarrow L$  is an isomorphism if  $A[f^{-1}] \rightarrow L[f^{-1}]$  is an isomorphism for all elements  $A[f^{-1}]$  of a covering family. Choose one which trivialises  $L$ . Then we are reduced to showing that every epimorphism of  $A$ -modules  $A \twoheadrightarrow A$  is an isomorphism. Morphisms of  $A$ -modules  $\phi : A \rightarrow A$  are all of the form  $a \mapsto ab$  for some  $b$ , namely,  $b = \phi(1)$ . Such a  $\phi$  is surjective if and only if 1 is in the image, i.e., if and only if  $ab = 1$  for some  $a$ , i.e., if and only if  $b$  is a unit. But in this case  $\phi$  is invertible with inverse  $a \mapsto ab^{-1}$ .

3. Suppose that  $I = \langle 0 \rangle$  is the zero ideal, but  $R$  is not the zero ring. Then there are no surjections  $A \otimes_R I \twoheadrightarrow L$  unless  $A$  is the zero ring. So in this case  $Bl_R I \cong \emptyset$  is the empty sheaf.

## 2 Strict transform

**Definition 4.** Let  $R \rightarrow R'$  be a ring homomorphism,  $I \subseteq R$  an ideal and set  $I' := IR'$ . Then there is a canonical morphism

$$Bl_{R'} I' \rightarrow Bl_R I$$

which sends a given  $R' \rightarrow A$  and  $[A \otimes_{R'} I' \twoheadrightarrow L]$  to the compositions  $R \rightarrow R' \rightarrow A$  and  $[A \otimes_R I \cong A \otimes'_{R'} R' \otimes_R I \twoheadrightarrow A \otimes_{R'} I' \twoheadrightarrow L]$ .

**Exercise 5.** Check that this is well defined and functorial. That is, check that:

1. if  $\text{Sym}_{R'}^n I' \rightarrow L$  factors through  $(I')^n$ , then  $\text{Sym}_R^n I \rightarrow L$  factors through  $I^n$ ,
2. given  $R' \rightarrow A \rightarrow B$ , the square

$$\begin{array}{ccc} Bl_{R'} I'(A) & \longrightarrow & Bl_R I(A) \\ \downarrow & & \downarrow \\ Bl_{R'} I'(B) & \longrightarrow & Bl_R I(B) \end{array}$$

commutes.

**Exercise 6.** Suppose that  $R \rightarrow R'$  is a flat ring homomorphism. I.e., the functor  $M \mapsto R' \otimes_R M$  preserves monomorphisms. Let  $I \subseteq R$  be an ideal, and set  $I' := IR'$ . Show that the canonical morphism

$$Bl_{R'} I' \rightarrow \text{Spec}(R') \times_{\text{Spec}(R)} Bl_R I$$

is an isomorphism.

**Example 7.** The morphism

$$Bl_{R'} I' \rightarrow \text{Spec}(R') \times_{\text{Spec}(R)} Bl_R I$$

is not always an isomorphism. For example, if  $R = \mathbb{Z}[x, y]$ ,  $I = \langle x, y \rangle$ , and  $R' = \mathbb{Z}$  (with  $R$ -algebra structure  $x, y \mapsto 0$ ) then  $I' = \langle 0 \rangle$  so  $Bl_{R'} I' = \emptyset$ . However, since  $I/I^2 \cong \mathbb{Z} \oplus \mathbb{Z}$  is free, the pullback along  $\text{Spec}(\mathbb{Z}) \rightarrow \text{Spec}(\mathbb{Z}[x, y])$  is  $\mathbb{P}^1$ , as discussed at the beginning of the lecture.

$$\begin{array}{ccccc} \emptyset \cong Bl_{R'} I' & \xrightarrow{\cong} & \mathbb{P}^1 \cong \text{Spec}(R') \times_{\text{Spec}(R)} Bl_R I & \longrightarrow & Bl_R I \\ & \searrow & \downarrow & & \downarrow \\ & & \text{Spec}(\mathbb{Z}) & \longrightarrow & \text{Spec}(\mathbb{Z}[x, y]) \end{array}$$

### 3 Descent for schemes

We want to define blowups along closed immersions of schemes. We will do this by a descent argument. We will use the following easy fact.

**Exercise 8.** Suppose that  $X$  is a sheaf and  $\{U_\lambda \rightarrow X\}_{\lambda \in \Lambda}$  is a family of open immersions such that each  $U_\lambda$  is a scheme and  $\sqcup U_\lambda \rightarrow X$  is an epimorphism of sheaves. Show that  $X$  is a scheme.

**Theorem 9.** Write  $\text{Sch}_R = \text{Sch}_{/\text{Spec}(R)}$  for the comma category. The assignment  $R \mapsto \text{Sch}_R$  is a 2-functor satisfying excision. That is, for every Zariski covering in  $\text{Aff}$  of the form  $\{U \rightarrow X, V \rightarrow X\}$ , the canonical functor

$$\Phi : \text{Sch}_{/X} \rightarrow 2\text{-lim} \left( \text{Sch}_{/U} \times \text{Sch}_{/V} \rightrightarrows \text{Sch}_{/W} \right) \quad (2)$$

is an equivalence of categories. Here  $W = U \times_X V$ .

*Proof.* To begin with we construct a left adjoint to (2). Suppose that we have morphisms of schemes  $Y_U \rightarrow U$ ,  $Y_V \rightarrow V$ ,  $Y_W \rightarrow W$  and an isomorphism of  $W$ -schemes  $Y_U \times_U W \xrightarrow{a} Y_W \xrightarrow{b} Y_V \times_V W$ . Define  $Y$  via the pushout of *sheaves*

$$\Psi(Y_V, Y_W, Y_U, a, b) := Y := Y_V \sqcup_{Y_W} Y_U.$$

This is clearly functorial in  $Y_\bullet$ . Since  $X = V \sqcup_W V$  it is equipped with a canonical morphism  $Y \rightarrow X$ . We show that  $Y$  is a scheme.

Since  $Y_V$  and  $Y_U$  are schemes, it suffices to show that  $Y_V, Y_U \rightarrow Y$  are open immersions and  $Y_V \sqcup Y_U \rightarrow Y$  is an epimorphism, see Exercise 8. Since  $Y = Y_V \sqcup_{Y_W} Y_U$ , the morphism  $Y_V \sqcup Y_U \rightarrow Y$  is automatically an epimorphism. For the open immersions, it suffices to show that the two squares

$$\begin{array}{ccc} Y_V & \longrightarrow & Y \\ \downarrow & & \downarrow \\ V & \longrightarrow & X \end{array} \quad \begin{array}{ccc} Y_U & \longrightarrow & Y \\ \downarrow & & \downarrow \\ U & \longrightarrow & X \end{array}$$

are cartesian. We have

$$V \times_X Y_W \cong V \times_X W \times_W Y_W \cong W \times_W Y_W \cong Y_W \quad (3)$$

and similarly,  $V \times_X Y_U \cong Y_W$  and  $V \times_X Y_V \cong Y_V$  so

$$V \times_X (Y_V \sqcup_{Y_W} Y_U) \cong (V \times_X Y_V \sqcup_{V \times_X Y_W} V \times_X Y_U) \cong (Y_V \sqcup_{Y_W} Y_W) \cong Y_V$$

So we have a functor  $\Psi$ . To conclude it suffices to show that the canonical natural transformations  $\Psi\Phi \rightarrow \text{id}$  and  $\text{id} \rightarrow \Phi\Psi$  are isomorphisms. That is, we want to show that for  $(Y_V, Y_U, Y_W, a, b)$  as above, and  $T \rightarrow X$  in  $\text{Sch}_{/X}$  the canonical morphisms

$$T \times_X U \sqcup_{T \times_X W} T \times_X V \rightarrow T$$

and

$$\begin{aligned} Y_U &\rightarrow Y \times_X U \\ Y_V &\rightarrow Y \times_X V \\ Y_W &\rightarrow Y \times_X W \end{aligned}$$

are isomorphisms. The first one is true because colimits are universal, and the second set of three is the same as Eq.(3).  $\square$

## 4 Blowups of schemes

**Corollary 10.** *Let  $Z \rightarrow X$  be a closed immersion of schemes. Then there exists an  $X$ -scheme*

$$Bl_X Z,$$

*unique up to unique isomorphism, such that for open immersion  $U := \text{Spec}(A) \rightarrow X$  we have an isomorphism*

$$U \times_X Bl_X Z \cong Bl_A I$$

*where  $\text{Spec}(A) \times_X Z = \text{Spec}(A/I)$ , and for every open immersion  $V = \text{Spec}(B) \rightarrow \text{Spec}(A) = U$  we have a commutative square*

$$\begin{array}{ccc} V \times_X Bl_X Z & \longrightarrow & Bl_B I B \\ \downarrow & & \downarrow (*) \\ V \times_U U \times_X Bl_X Z & \longrightarrow & V \times_U Bl_A I \end{array}$$

*where  $(*)$  is the canonical morphism from Definition 4.*

*Proof.* Since  $X = \text{colim}_{\text{Aff}/X}^{\text{open}} j \text{Spec}(A)$ , by Theorem 9 (and the 2-category version of the excision theorem from last week) we have

$$\text{Sch}_X \xrightarrow{\sim} 2\text{-lim}_{\substack{\text{Spec}(A) \rightarrow X \\ \text{open}}} \text{Sch}_A.$$

Above we have constructed:

1. For every  $A$  an  $A$ -scheme  $Bl_A I$ ,
2. For every open immersion  $\text{Spec}(B) \rightarrow \text{Spec}(A)$  an isomorphism  $Bl_B I B \xrightarrow{\sim} \text{Spec}(B) \times_{\text{Spec}(A)} Bl_A I$ .

So to finish defining an object of  $2\text{-lim} \text{Sch}_A$  we have to show that for open immersions  $W \rightarrow V \rightarrow U \rightarrow X$  with  $W, V, U$  affine, the square (below) of isomorphisms is commutative. Here we write  $Z_U, Z_V, Z_W$  for the respective pullbacks of  $Z \rightarrow X$  to  $U, V, W$ . For this it suffices to check that the triangle on the right is commutative.

$$\begin{array}{ccc} Bl_W Z_W & \longrightarrow & W \times_V Bl_V Z_V \\ \downarrow & & \downarrow \\ W \times_V V \times_U Bl_U Z_U & \longrightarrow & W \times_U Bl_U Z_U \end{array} \qquad \begin{array}{ccc} Bl_W Z_W & \longrightarrow & Bl_V Z_V \\ & \searrow & \downarrow \\ & & Bl_U Z_U \end{array}$$

This follows from the definition, Def.4.  $\square$

## 5 Deformation to the normal cone

Suppose that  $\text{Spec}(B) = \text{Spec}(A/I) = Z \rightarrow X = \text{Spec}(A)$  is a closed immersion of affine schemes. Suppose furthermore, that  $I/I^2 \cong B^{\oplus c}$  for some  $c$ . If  $U = \text{im}(\sqcup_{a \in I} \text{Spec}(A[a^{-1}]) \rightarrow \text{Spec}(A))$  is the open complement to  $Z$ , as discussed in Example 3, we obtain the following cartesian squares

$$\begin{array}{ccccc} U & \longrightarrow & \text{Bl}_X Z & \longleftarrow & Z \times \mathbb{P}^{c-1} \\ \parallel & & \downarrow & & \downarrow \\ U & \longrightarrow & X & \longleftarrow & Z \end{array}$$

Now we consider  $X \times \mathbb{A}^1 = \text{Spec}(A[t])$  with closed subscheme  $Z \times \{0\} = \text{Spec}(A[t]/J)$  where  $J = IA[t] + tA[t]$ . One can check that we have  $J/J^2 \cong (I/I^2) \oplus B$ . So this leads to the cartesian square

$$\begin{array}{ccc} \text{Bl}_{X \times \mathbb{A}^1} Z \times \{0\} & \longleftarrow & Z \times \mathbb{P}^c \\ \downarrow & & \downarrow \\ X \times \mathbb{A}^1 & \longleftarrow & Z \times \{0\} \end{array}$$

Moreover, we have the zero section  $s_0 : X \rightarrow X \times \mathbb{A}^1$  and it's induced square

$$\begin{array}{ccc} \text{Bl}_X Z & \longrightarrow & \text{Bl}_{X \times \mathbb{A}^1} Z \times \{0\} \\ \downarrow & & \downarrow \\ X & \longrightarrow & X \times \mathbb{A}^1 \end{array}$$

This square is *not* cartesian, but the horizontal morphisms are none-the-less closed immersions. We want to consider the open complement

$$\mathcal{D}_{X,Z} := \text{Bl}_{X \times \mathbb{A}^1} Z \times \{0\} \setminus \text{Bl}_X Z$$

with it's canonical morphism

$$\pi : \mathcal{D}_{X,Z} \rightarrow \mathbb{A}^1.$$

More precisely,  $\mathcal{D}_{X,Z}$  is the union of all open immersions  $V \subseteq \text{Bl}_{X \times \mathbb{A}^1} Z \times \{0\}$  whose intersection with  $\text{Bl}_X Z$  is empty.

This comes equipped with the following closed immersion  $Z \times \mathbb{A}^1 \rightarrow \mathcal{D}_{X,Z}$ . Since the ideal  $tB[t] \subseteq B[t]$  of  $Z \times \{0\} \rightarrow Z \times \mathbb{A}^1$  is generated by a single non-zero divisor, there is an isomorphism  $tB[t] \cong B[t]$ , so  $\text{Bl}_{Z \times \mathbb{A}^1} Z \times \{0\} \cong Z \times \mathbb{A}^1$  and we get a canonical morphism

$$Z \times \mathbb{A}^1 \cong \text{Bl}_{Z \times \mathbb{A}^1} Z \times \{0\} \rightarrow \text{Bl}_{X \times \mathbb{A}^1} Z \times \{0\}.$$

One can check that this doesn't intersect  $Z \times \mathbb{P}^{c-1} \subseteq Z \times \mathbb{P}^c$  so it factors as

$$\iota : Z \times \mathbb{A}^1 \rightarrow \mathcal{D}_{X,Z}.$$

By construction, away from  $\{0\} \subseteq \mathbb{A}^1$  the maps  $\mathcal{D}_{X,Z} \rightarrow Bl_{X \times \mathbb{A}^1} Z \times \{0\} \rightarrow X$  are isomorphisms. That is, if we pullback  $\iota$  along  $\mathbb{A}^1 \setminus \{0\}$  we get the closed immersion  $Z \times (\mathbb{A}^1 \setminus \{0\}) \rightarrow X \times (\mathbb{A}^1 \setminus \{0\})$ . On the other hand, by construction, over  $\{0\}$  we have a closed immersion of the form  $Z \rightarrow Z \times \mathbb{A}^c$  and one can check that this is the zero section. So we have the following cartesian squares.

$$\begin{array}{ccccc}
Z \times (\mathbb{A}^1 \setminus 0) & \longrightarrow & Z \times \mathbb{A}^1 & \longleftarrow & Z \times \{0\} \\
\downarrow & & \downarrow & & \downarrow \\
X \times (\mathbb{A}^1 \setminus 0) & \longrightarrow & \mathcal{D}_{X,Z} & \longleftarrow & Z \times \mathbb{A}^c \\
\downarrow & & \downarrow & & \downarrow \\
\mathbb{A}^1 \setminus 0 & \longrightarrow & \mathbb{A}^1 & \longleftarrow & \{0\}
\end{array}$$

That is, we have “deformed” the closed immersion  $Z \rightarrow X$  into the zero section  $Z \rightarrow Z \times \mathbb{A}^c$  of a vector bundle. In fact, all of this works with the weaker assumption that  $I/I^2$  is finite rank projective, but not necessarily free, in which case we get the vector bundle  $\text{Spec}(\oplus \text{Sym}_B^d I/I^2)$  over  $Z = \text{Spec}(B)$ .

## 6 Further exercises

**Exercise 11** ((Harder) The standard open affine covering). Consider the closed immersion  $Bl_R I \rightarrow \mathbb{P}_R^\Lambda$  from Proposition 2. Recall from the Schemes lecture that for each  $\lambda \in \Lambda$  we get an open immersion  $\mathbb{A}_R^{\Lambda \setminus \{\lambda\}} \rightarrow \mathbb{P}_R^\Lambda$ , and these form an open affine covering of  $\mathbb{P}_R^\Lambda$ . This induces an open affine covering

$$\sqcup_{\lambda \in \Lambda} Bl_R I \times_{\mathbb{P}_R^\Lambda} \mathbb{A}_R^{\Lambda \setminus \{\lambda\}} \rightarrow Bl_R I.$$

Show that  $Bl_R I \times_{\mathbb{P}_R^\Lambda} \mathbb{A}_R^{\Lambda \setminus \{\lambda\}}$  is the affine scheme associated to the ring

$$S_{r_\lambda} := \text{colim}(R \xrightarrow{r_\lambda} I \xrightarrow{r_\lambda} I^2 \xrightarrow{r_\lambda} \dots)$$

where the transition morphisms are multiplication by  $r_\lambda$ . Note that  $S_{r_\lambda}$  can be considered as the subring

$$\left\{ \frac{m}{r_\lambda^n} \in R[r_\lambda^{-1}] \mid m \in I^n \right\}$$

by assembling the morphisms  $I^n \rightarrow R[r_\lambda^{-1}]; m \mapsto \frac{m}{r_\lambda^n}$ .

**Definition 12.** As with  $\mathbb{P}^n$ , the assignment

$$\mathcal{O}(1) : [A \otimes_R I \twoheadrightarrow L] \mapsto L$$

(where we make a choice in each equivalence class) defines an invertible module on the presheaf  $Bl_R I$ .



**Exercise 13** (The invertible module  $\mathcal{O}(1)$ ). Consider the assignment

$$\mathcal{O}(1) : [A \otimes_R I \twoheadrightarrow L] \mapsto L$$

(where we make a choice in each equivalence class).

Show that  $\mathcal{O}(1)$  actually is a quasi-coherent module. That is, show that one can choose an isomorphism  $\phi_{s,f} : B \otimes_A \mathcal{O}(1)(s) \xrightarrow{\sim} \mathcal{O}(1)(sf)$  for every pair of morphisms  $f : \text{Spec}(B) \rightarrow \text{Spec}(A)$  and  $s \in \text{Bl}_R I(A) \cong \text{hom}(\text{Spec}(A), \text{Bl}_R I)$ , such that given a third morphism  $g : \text{Spec}(C) \rightarrow \text{Spec}(B)$  the square

$$\begin{array}{ccc} C \otimes_B B \otimes_A \mathcal{O}(1)(s) & \longrightarrow & C \otimes_A \mathcal{O}(1)(s) \\ C \otimes_B \phi_{s,f} \downarrow & & \downarrow \phi_{s,fg} \\ C \otimes_B \mathcal{O}(1)(sf) & \xrightarrow{\phi_{sf,g}} & \mathcal{O}(1)(sfg) \end{array}$$

commutes.

**Example 14.** Show that  $\text{Bl}_{\mathbb{Z}[x_1, \dots, x_n]} \langle x_1, \dots, x_n \rangle$  admits an open affine covering by the schemes  $\text{Spec} \mathbb{Z}[\frac{x_1}{x_i}, \dots, \frac{x_{i-1}}{x_i}, x_i, \frac{x_{i+1}}{x_i}, \dots, \frac{x_n}{x_i}] \cong \mathbb{A}^n$ .